AIN SHAMS UNIVERSITY

Faculty of Science Mathematics Department

THEORY AND APPLICATION OF ZERO-KNOWLEDGE **PROOFS**



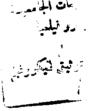
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LIST OF ABBREVIATIONS

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p	The complexity class of average polynomial time problems. The class of languages recognized by an Arthur-Merlin
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	The class of languages recognized by all Artiflat-World game of $q(n)$ message exchanged.
BCS* E	Bit commitment scheme. Bit commitment scheme without any intractability assumptions. The i^{th} bit of a pseudo-random sequence on a seed s .
CZK	The class of problems whose complements are in <i>NP</i> . The class of languages which have computational zero-knowledge proofs.
GI GNP G GSC GMW G $G_j(s)$	Deterministic Turing machine. Graph isomorphism problem. Graph nonisomorphism problem. Graph 3-colourability. Goldreich, Micali and Wigderson. The first j bits of a pseudo-random sequence on a seed $s \in \{0, 1\}^n$.
IP .	Interactive proof system. The class of languages which have interactive proofs. Interactive Turing machine.
	The class of languages which have multi-prover interactive proofs. The simulator of V^* .
	The class of languages recognizable by nondeterministic polynomial-time Turing machine.
PZK	The probability that (P, V) accepts the common input x . The class of languages which have perfect zero-knowledg proofs. The conversation space between P and V , on input x .
	k-prover interactive protocol.

PSPACE	The class of all languages recognizable by polynomial space bounded <i>DTM</i> programs that halts on all inputs.
QRA QBF	Quadratic residuosity assumption. Quantified Boolean formula problem.
RNP RTP RSA	The class of randomized decision problems. Randomized Tiling problem. Riverst, Shamir and Adleman.
$x \oplus y$ $x \in \{0, 1\}^n$ SZK Sym(N)	The bit by bit exclusive-or of bit strings x and y . x is a string of n bits. The class of languages which have statistical zero-knowledge proofs. The set of all possible permutations on N where $N = \{1, 2,, n\}$.
V V* V _P (x) View _{(P1,, Pk, 17} (x)	Verifier. Cheating verifier. V 's output after interacting with P on a common input x . The verifier's view during the protocol.
$ZKIPS$ Z_p	Zero-knowledge interactive proof systems. The set of integers $\{0,, p-1\}$, where p is a prime. We can view Z_p as a group with respect to addition modulo p .
Z* _p	The set of integers $\{z \in \{0,, p-1\}: gcd(z, p) = 1\}$. We can view Z^*_p as a group with respect to multiplication modulo p .
3SAT μ ν(n)	3-Satisfiability problem. Probability distribution function. Any function vanishing faster than the inverse of any
[] {0, 1} ⁿ "	polynomial in n . Decision problem. The set of all bit strings of length n . The concatenation of n l bits.



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